

# Convoy Effect Elimination in FCFS Scheduling



# Sambath M., K. Padmaveni, Linda Joseph, Ravi S., J. Thangakumar, D. John Aravindhar

Abstract: One of the important activities of operating systems is process scheduling. There are many algorithms available for scheduling like First Come First Served, Shortest Job First, Priority Scheduling and Round Robin. The fundamental algorithm is First Come First Served. It has some drawback of convoy effect. Convoy effect occurs when the small processes are waiting for lengthy process to complete. In this paper novel method is proposed to reduce convoy effect and to make the Scheduling optimal which reduces average waiting time and turnaround time.

Keywords: convoy effect, scheduler, throughput, turnaround time.

### I. INTRODUCTION

CPU Scheduling is the process of assigning the process to the processor for execution. In memory multiple programs are stored and ready for execution. The performance of the processor depends on how the programs are dispatched to the processor for execution. In First Come First Served (FCFS) algorithm the processes are dispatched in the order of arrival to the processor. Scheduling is done by the scheduler[1]. The scheduling algorithm reflects the overall performance of the processor. There are three types of schedulers: [1]

- 1. Long-term scheduler (Job scheduler): It is responsible for loading the programs from the memory into the ready queue.
- 2. Short-term scheduler (CPU Scheduler): Schedules the selected process for execution and gives control to the CPU[2].
- 3. Medium-term scheduler: Time sharing systems uses this kind of scheduler. The process is preempted for some reason and resumed back after sometime. [9]

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- The performance of the scheduling algorithm is estimated based on various criteria like CPU utilization, Throughput, Turnaround time, waiting time and response time.[3]
- 1. CPU utilization indicates how effectively CPU is utilized.[8]
- 2. Throughput denotes the number of processes completed per unit time.[4]
- 3. Turnaround Time is the time needed for the process to complete its execution.[5]
- 4. Waiting Time is the amount of time a process waiting in the queue.
- 5. Response Time is the time between submission of request for execution and till the first response is produced not the output.

#### II. FCFS SCHEDULING

The processes are loaded from memory into ready queue. From the ready queue the processes are taken based on their arrival time and assigned to the processor for execution[6] as shown in Figure 1.



#### Figure 1 FCFS Scheduling

The table 1 lists the processes and their execution time (burst time). The FCFS scheduling (Figure 2) is applied and the average waiting time and turnaround time are calculated and it is listed in table 2.

| e |
|---|
| l |

| S.No. | Process | Burst Time |
|-------|---------|------------|
| 1     | P1      | 15         |
| 2     | P2      | 19         |
| 3     | P3      | 35         |
| 4     | P4      | 40         |
| 5     | P5      | 50         |
| 6     | P6      | 58         |
| 7     | P7      | 3          |
| 8     | P8      | 5          |
| 9     | P9      | 7          |
| 10    | P10     | 6          |



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# **Convoy Effect Elimination in FCFS Scheduling**

|     | 1    |       | 1          | 1            | 1           |             | 1        | 1        |         |
|-----|------|-------|------------|--------------|-------------|-------------|----------|----------|---------|
| P1  | P2   | P3    | P4         | P5           | P6          | P7          | P8       | P9       | P10     |
| 0   | 15   | 34    | 69         | 109          | 159         | 217         | 220      | 225      | 232     |
|     |      |       | Б          | Source 2 Con | att Chant w | ing ECES    |          |          | 238     |
|     |      | _     | г<br>      | igure 2 Gai  |             | sing FCF5   |          |          |         |
| r   |      |       | Table 2 Wa | iting time a | nd turnarou | ind time us | ing FCFS |          |         |
| Pro | cess | Burst | Time       |              | Waiting t   | ime         |          | Turnarou | nd time |
| P1  |      | 15    |            | 0            |             |             | 15       |          |         |
| P2  |      | 19    |            | 15           |             |             | 34       |          |         |
| Р   | 3    | 35    |            |              | 34          | 4 69        |          |          |         |
| Р   | 4    | 40    |            | 40 69        |             | 109         |          | )        |         |
| P5  |      | 50    |            |              | 109         |             |          | 159      | )       |
| Р   | P6   |       | 58         |              | 159         |             |          | 217      | 7       |
| P7  |      | 3     |            |              | 217         |             |          | 220      | )       |
| P8  |      | 5     |            |              | 220         |             |          | 225      | 5       |
| Р   | P9   |       | 7          |              | 225         |             |          | 232      | 2       |
| P   | 10   | 6     |            |              | 232         |             |          | 238      | 3       |
|     |      | Ave   | rage       |              | 128         |             |          | 151.     | .8      |

The average waiting time is 128 unit time and average turnaround time is 151.8 unit time. From the table 2 it is observed that the short processes are available after the lengthy processes[7]. Due to this small processes are waiting till the completion of lengthy processes. This leads to increase in waiting time as well as turnaround time. This effect is called as convoy effect. To reduce this effect novel approach – BinFCFS algorithm is proposed.

# **III. BINFCFS ALGORITHM**

- 1. Find out the minimum and maximum in burst time.
- 2. Create the number of bins depends on the range of burst time.
- 3. Group the processes and put it in the bins based on the range of burst time.

| Bin 0<br>(0-9) | Bin 1<br>(10-19) | Bin 2<br>(20-29) |
|----------------|------------------|------------------|
| P7             | P1               |                  |
| P8             | P2               |                  |
| P9             |                  |                  |
| P10            |                  |                  |



4. Take the processes from the bins in the order and execute.

For the above scenario, the processes are arranged into the bins based on their burst time as shown in Figure 3. Start from Bin 0 executes P7, P8, P9 and P10. Then continue with Bin1, Bin2 and so on.

For the above scenario, the processes are arranged into the bins based on their burst time as shown in Figure 3. Start from Bin 0 executes P7, P8, P9 and P10. Then continue with Bin1, Bin2 and so on.

The waiting time and turnaround time for each process is calculated using the proposed method -BinFCFS algorithm(Figure 4) and is given in the table 3.

| Bin 3<br>(30-39) | Bin 4<br>(40-49) | Bin 5<br>(50-59) |
|------------------|------------------|------------------|
| P3               | P4               | P5               |
|                  |                  | P6               |
|                  |                  |                  |
|                  |                  |                  |

Figure 3 Bins holding the processes



Figure 4 Gantt Chart using BinFCFS algorithm



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| Process | Burst Time | waiting time | Turnaround time |
|---------|------------|--------------|-----------------|
| P7      | 3          | 0            | 3               |
| P8      | 5          | 3            | 8               |
| P9      | 7          | 8            | 15              |
| P10     | 6          | 15           | 21              |
| P1      | 15         | 21           | 36              |
| P2      | 19         | 36           | 55              |
| P3      | 35         | 55           | 90              |
| P4      | 40         | 90           | 130             |
| P5      | 50         | 130          | 180             |
| P6      | 58         | 180          | 238             |
|         | Average    | 85.33        | 121.5           |

# Table 3 Waiting time and turnaround time using BinFCFS

# IV. RESULT AND DISCUSSION

The average waiting time and turnaround time for the above problem using FCFS algorithm is 128 and 151.8 as given in table 4. It is reduced to 85.33 and 121.5 respectively using the proposed BinFCFS algorithm.



## **Table 4 Performance Comparison**

| Algorithm            | Average<br>waiting time | Average<br>Turnaround time |
|----------------------|-------------------------|----------------------------|
| FCFS                 | 128                     | 151.8                      |
| BinFCFS<br>Algorithm | 85.33                   | 121.5                      |



# Figure 5 a) Performance comparison for Average waiting time b) Performance comparison of average turnaround time

## V. CONCLUSION

The FCFS algorithm schedules the processes in the order of arrival irrespective of whether it is a short process or long process. Due to this convoy effect is occurred. This effect is avoided by the proposed BinFCFS algorithm by grouping the short processes and executing them in the first.

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